



# Polynomial interpolation and counting lattice points in polytopes

I have polytope (dimension d)  
and I want :

Compute volume of  
this polytope  
OR  
Compute lattice points  
in this polytope

Compute polynomial of degree d

Algorithm?

Barvinok

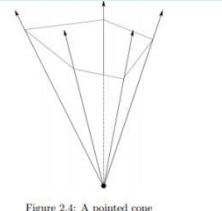


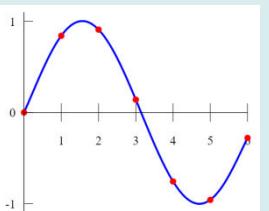
Figure 2.4: A pointed cone

The most common method  
Use power series

Clauss method is based  
on the polynomial  
interpolation :

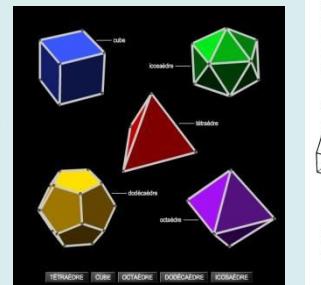
For polytope of  
dimension d we need to  
compute **d+1 values**

Clauss

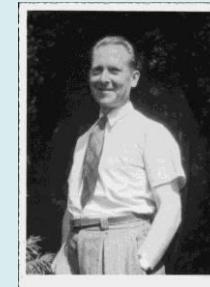


Our method is based on the  
polynomial interpolation :  
For polytope of dimension d  
**1 value** or just an  
**approximate** of this value  
is sufficient to generate  
(Ehrhart) polynomial

Convexe      Concave



Ehrhart theorem



$L_p(n) = a_d n^d + a_{d-1} n^{d-1} + a_{d-2} n^{d-2} + \dots + a_1 n^1 + a_0 ; a_i = c_i / d!$  avec  $c_i \in \mathbb{Z}$

$$L_p(t) = 2t + L'_p(t) = 120 + 1452t + 6120t^2 + 13200t^3 + 12720t^4 + 4608t^5 + t^6$$

$$L_p(10000002) = 460858800 + 10^9$$

1 value to compute  $L(t)$

$$c_5 = 10000002 \cdot \frac{L_p(10000002)}{10000002^5} = 4608$$

$$c_4 = 10000002 \cdot \frac{L_p'(10000002)}{10000002^4} = 12720$$

$$c_3 = 10000002 \cdot \frac{L_p''(10000002)}{10000002^3} - c_2 \cdot 10000002^2 - c_1 \cdot 10000002^1 + 12200$$

$$c_2 = 10000002 \cdot \frac{L_p'''(10000002)}{10000002^2} - c_1 \cdot 10000002^1 - c_0 \cdot 10000002^0 = 6420$$

$$c_1 = 10000002 \cdot \frac{L_p''''(10000002)}{10000002^1} - c_0 \cdot 10000002^0 = 1452$$

$$c_0 = 10000002$$

First point

Lower and upper bound on the coefficients  
of Ehrhart polynomial

• Upper bound Ulrich Betke et Peter McMullen (1985)

• Lower bound Martin Henk et Makoto Tagami (2009)

$$c_r < d! \left( (-1)^{d-r} \text{stirl}(d, r) \text{vol}(P) + (-1)^{d-r-1} \frac{\text{stirl}(d, r+1)}{(d-1)!} \right) = \alpha_r$$

$$c_r > d! \left( \frac{1}{d!} \left( (-1)^{d-r} \text{stirl}(d+1, r+1) + (d! \text{vol}(P) - 1) M_{r,d} \right) \right) = \beta_r$$



$$\frac{L_p(n)}{n^d} = c_d + \underbrace{\frac{c_{d-1}n^{d-1} + c_{d-2}n^{d-2} + \dots + c_1n^1 + c_0}{n^d}}_{\substack{\text{Lim} \rightarrow 0 \\ n \rightarrow \infty}}$$

$$\left[ \frac{L_p(k)}{k^d} \right] = c_d + \left[ \frac{c_{d-1}k^{d-1} + c_{d-2}k^{d-2} + \dots + c_1k^1 + c_0}{k^d} \right] = c_d = \text{vol}(P)$$

Work in progress.

(This example : 2 variable. General formula for p variables looks like this.)

Theorem 6.1. Let  $(\beta_r, c_r, \alpha_r) \in \mathbb{Z}^3$  as  $\beta_r \leq c_r \leq \alpha_r$ .  
Let  $P(n, m)$  an integer parametric polynomial :

$$P(n, m) = \sum_{r=0}^p \sum_{j=0}^q c_{r,j} n^r m^j = \sum_{r=0}^p g_r(m) n^r$$

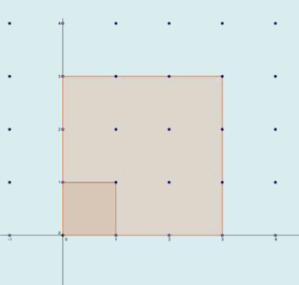
$$P(n, m) = (\sum_{j=0}^q c_{p,j} m^j) n^p + (\sum_{j=0}^q c_{p-1,j} m^j) n^{p-1} + \dots + (\sum_{j=0}^q c_{0,j} m^j)$$

We can have equivalent to lemma 1.1 and 1.2 for parametric polynomial so :

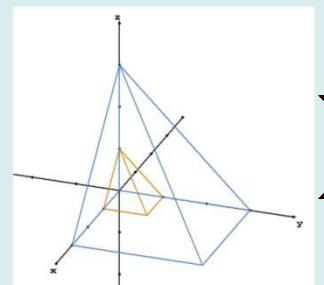
- $g_p(m) = \left\lfloor \frac{P(k_{max}, m)}{k_{max}^p} \right\rfloor$
- $g_r(m) = \left\lfloor \frac{P(k_{max}, m) - \sum_{r=0}^{p-1} g_r(m) k_{max}^r}{k_{max}^r} \right\rfloor \forall r \in [0, p-1]$
- $c_{r,q} = \left\lfloor \frac{g_q(k_{max}) - \sum_{j=0}^{q-1} c_{r,j} k_{max}^j}{k_{max}^q} \right\rfloor$
- $c_{r,j} = \left\lfloor \frac{g_j(k_{max}) - \sum_{i=j+1}^q c_{r,i} k_{max}^i}{k_{max}^j} \right\rfloor \forall j \in [0, q-1]$

Integer dilates of polytopes

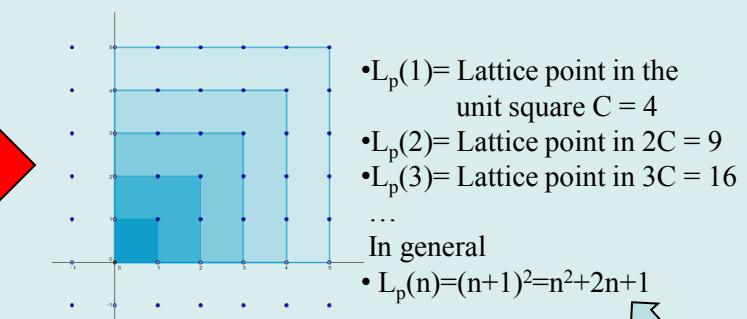
Square (2D)



Simplex (3D)



Lattice point in dilating unit square



- $L_p(1)$  = Lattice point in the unit square  $C = 4$
- $L_p(2)$  = Lattice point in  $2C = 9$
- $L_p(3)$  = Lattice point in  $3C = 16$

In general  

$$L_p(n) = (n+1)^2 = n^2 + 2n + 1$$

- Polynomial interpolation :  
Need to compute a,b,c
- $an^2 + bn + c = 4$
  - $an^2 + bn + c = 9$
  - $an^2 + bn + c = 16$

General formula

**Lemma 1.1.** Let  $\alpha_r \in \mathbb{Z}$ , we define  $f_{\alpha,d-1}$  as :

$$f_{\alpha,d-1}(t) = \frac{\sum_{r=0}^{d-1} \alpha_r t^r}{t^d}$$

So :

- $\lim_{t \rightarrow \infty} f_{\alpha,d-1}(t) = 0$
- $\exists k_{\alpha,d-1} \in \mathbb{N}$  as  $\forall t > k_{\alpha,d}$  then  $|f_{\alpha,d-1}(t)| < \frac{1}{2}$

**Lemma 1.2.** Let  $(\beta_r, c_r, \alpha_r) \in \mathbb{Z}^3$  as  $\beta_r \leq c_r \leq \alpha_r$ . Let  $f_{\alpha,d}$ ,  $f_{\beta,d}$  et  $f_{c,d}$  function define in lemma 1.1. we define  $k_{\max,d-1}$ ,  $k_{\max}$  as :

- $k_{\max,d-1} = \text{Max}\{k_{\alpha,d-1}, k_{\beta,d-1}\}$
- $k_{\max} = \text{Max}\{k_{\max,r} \forall r \in [0, d-1]\}$

So  $\forall t > k_{\max}$ ,  $\forall r \in [0, d-1]$  we have :

$$|f_{c,r}(t)| = 0$$

Maximum margin of error  
of our approximation. We can control this upper bound. If we add another approximation this bound increase.

**Theorem 1.3.** Let  $(\beta_r, c_r, \alpha_r) \in \mathbb{Z}^3$  as  $\beta_r \leq c_r \leq \alpha_r$ . Let  $P(t)$  an integer polynomial :

$$P(t) = \sum_{r=0}^d c_r t^r$$

So :

Combining lemma 1.1 and 1.2 we have :

- $c_d = \left\lfloor \frac{P(k_{\max})}{k_{\max}^d} \right\rfloor$
- $c_r = \left\lfloor \frac{P(k_{\max}) - \sum_{i=r+1}^d c_i k_{\max}^i}{k_{\max}^r} \right\rfloor \forall r \in [0, d-1]$

Starting from only one estimate  $P(k_{\max})$ , we can generate the Ehrhart polynomial.

Starting from an approximation of  $P(k_{\max})$ , we can generate the Ehrhart polynomial.

**Parametric polytope** : We need just one value (or approximation) to generate any parametric multivariate polynomials.